Distance vector protocols

Distance Vector Routing Principles

Routing loops and countermeasures to loops

Bellman-Ford route calculations

RIP

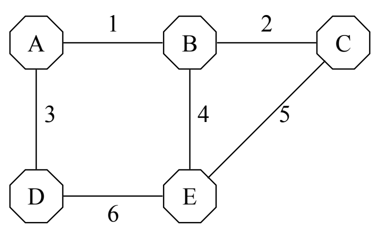
Distance Vector Routing Principles

RIP - Routing Information Protocol is a basic protocol for interior routing

- RIP is a distance vector protocol
  - Based on the Bellman-Ford algorithm
- Routing table contains information about other known nodes
  - distance
  - link
- The nodes periodically send distance vectors based on the routing tables on all their links
- The nodes update their routing table with received distance vectors

Let us study DV protocol principles

Example network with nodes A, B, C, D, E and links 1, 2, 3, 4, 5, 6.



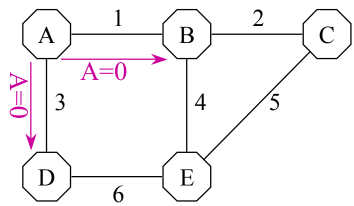
Initial state: Nodes know their own addresses and interfaces, nothing more.

Node A creates its routing table:

From A to node	Link	Cost
A	local	0

The corresponding DV is: A=0

Generation of routing tables starts when all routers send their DVs on all interfaces



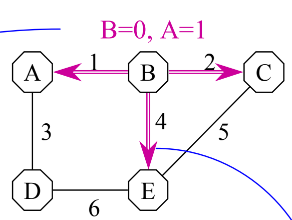
Let's look at reception in Node B:

From node B to	Link	Cost
B	local	0

1. B receives the DV A=0
2. B increments the DV + 1 => A=1 and
3. B looks for the result in its routing table, no match
4. B adds the result to its RT, result is

From node B to	Link	Cost
B	local	0
A	1	1

B creates its own DV and sends it to all neighbors



C to	Link	Cost
C	-	0
B	2	1
A	2	2

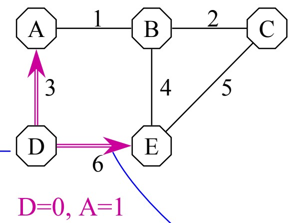
A to	Link	Cost
A	-	0
B	1	1

A=2 > A=0

E to	Link	Cost
E	-	0
B	4	1
A	4	2

D sends its distance vector to all neighbors

A to	Link	Cost
A	-	0
B	1	1
D	3	1

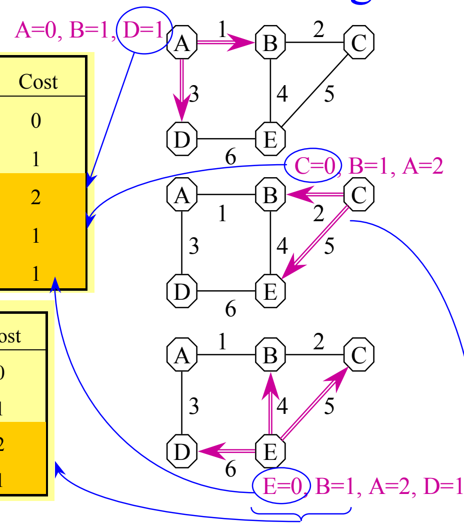


C to	Link	Cost
C	-	0
B	2	1
A	2	2

E to	Link	Cost
E	-	0
B	4	1
A	4	2
D	6	1

A=2 == A=2 ⇒ no change

Nodes with changed RTs create DVs and send them to neighbors



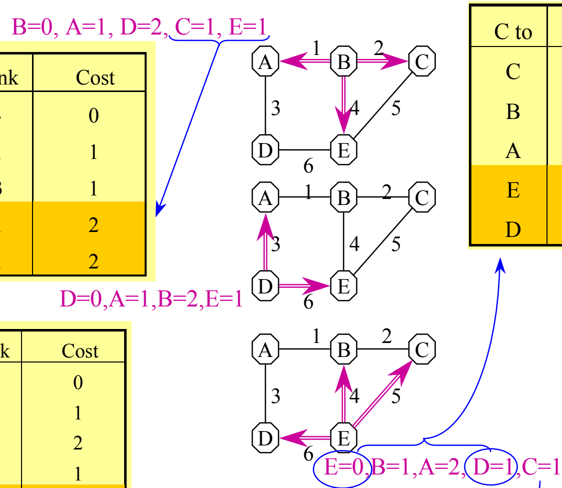
B to	Link	Cost
B	-	0
A	1	1
D	1	2
C	2	1
E	4	1

D to	Link	Cost
D	-	0
A	3	1
B	3	2
E	6	1

E to	Link	Cost
E	-	0
B	4	1
A	4	2
D	6	1
C	5	1

Again the changes are sent ...

A to	Link	Cost
A	-	0
B	1	1
D	3	1
C	1	2
E	1	2

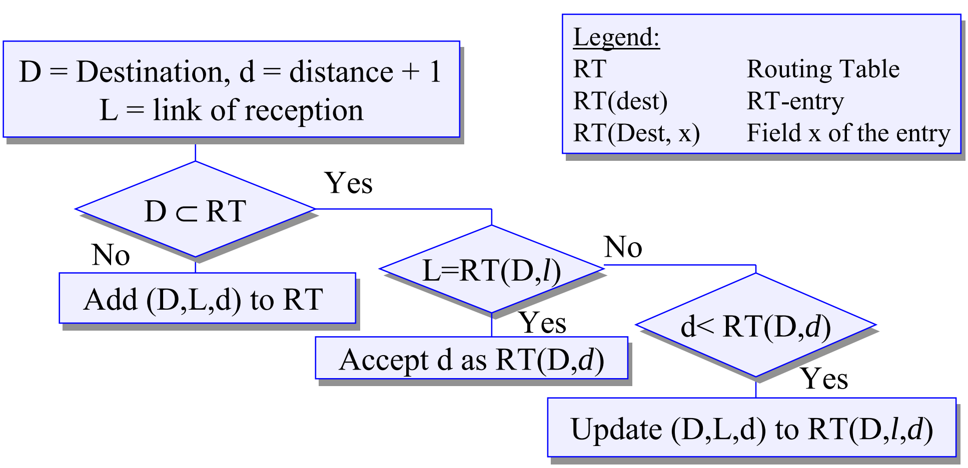


C to	Link	Cost
C	-	0
B	2	1
A	2	2
E	5	1
D	5	2

D to	Link	Cost
D	-	0
A	3	1
B	3	2
E	6	1
C	6	2

A, D, and C create new DVs, send them, but they have no impact.

Processing of Received Distance Vectors



Note: this is simplified, shows only the principle!

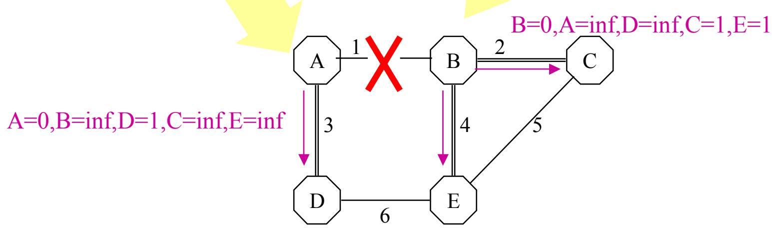
A link breaks...

A round of updates starts on link failure

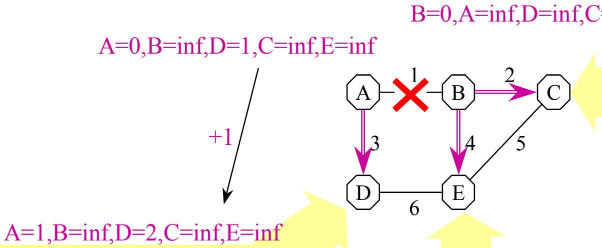
A:to	Link	Cost
A	-	0
B	1	inf.
D	3	1
C	1	inf.
E	1	inf.

A gives an infinite distance to the nodes reached through link 1

B to	Link	Cost
B	-	0
A	1	inf
D	1	inf
C	2	1
E	4	1



D, E ja C update their Routing Tables



D to	Link	Cost
D	-	0
A	3	1
B	3	inf
E	6	1
C	6	2

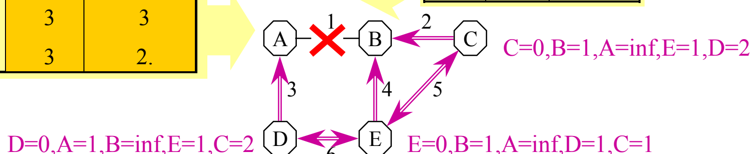
C to	Link	Cost
C	-	0
B	2	1
A	2	inf
E	5	1
D	5	2

E to	Link	Cost
E	-	0
B	4	1
A	4	inf
D	6	1
C	5	1

D, C, E generate Distance Vectors...

A to	Link	Cost
A	-	0
B	1	inf.
D	3	1
C	3	3
E	3	2.

B to	Link	Cost
B	-	0
A	1	inf
D	4	2
C	2	1
E	4	1



D to	Link	Cost
D	-	0
A	3	1
B	6	2
E	6	1
C	6	2

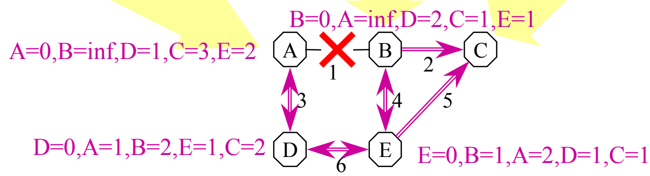
E to	Link	Cost
E	-	0
B	4	1
A	6	2
D	6	1
C	5	1

A, B, D, E generate their Distance Vectors

A to	Link	Cost
A	-	0
B	3	3
D	3	1
C	3	3
E	3	2.

B to	Link	Cost
B	-	0
A	4	3
D	4	2
C	2	1
E	4	1

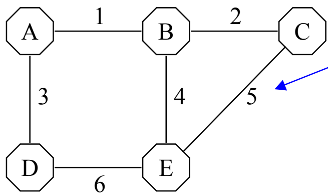
C to	Link	Cost
C	-	0
B	2	1
A	5	3
E	5	1
D	5	2



The result is that all nodes are able to communicate with all other nodes.

Routing loops

DV -protocol may create a transient routing loop

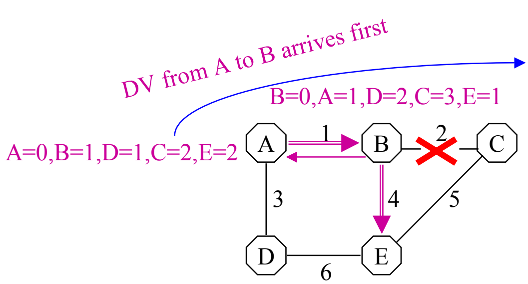


Let's assume that cost of loop 5 is 8. A stable initial state for routes to C would be:

x to C	Link from x	Cost
A->C	1	2
B->C	2	1
C->C	-	0
D->C	3	3
E->C	4	2

Let's just look at the first link of each route.

Link 2 fails



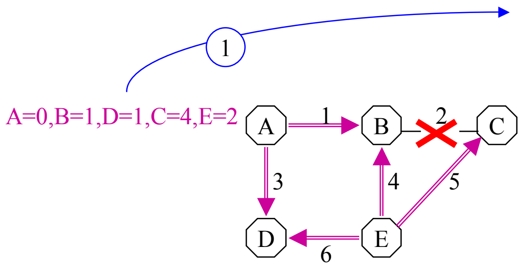
Intermediate state

x to C	Link from x	Cost
A->C	1	2
B->C	2	inf
C->C	-	0
D->C	3	3
E->C	4	2

x to C	Link from x	Cost
A->C	1	4
B->C	1	3
C->C	-	0
D->C	3	3
E->C	4	4

All packets to C are sent to B. B sends them to A. A sends them back to B… until TTL=0. (Bouncing effect)

A and E send their Distance Vectors



x to C	Link rom x	Cost
A->C	1	4
B->C	1	5
C->C	-	0
D->C	3	5
E->C	4	4

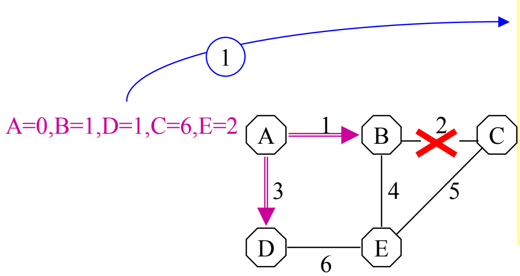
2

B generates a new DV:
B=0,A=1,D=2,C=5,E=1

⇒ Distance seen by A to C grows to 6

Distance vectors sent by C do not change anything because of high link cost

A sends a new DV



x to C	Link from x	Cost
A->C	1	6
B->C	1	7
C->C	-	0
D->C	3	7
E->C	4	6

2

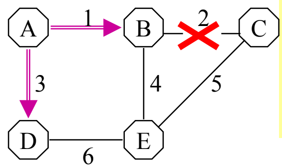
B generates a new DV
B=0,A=1,D=2,C=7,E=1

⇒ Distance seen by A to C grows to 8

A sends a new DV

1

A=0,B=1,D=1,C=8,E=2



x to C	Link from x	Cost
A->C	1	8
B->C	1	9
C->C	-	0
D->C	3	9
E->C	4	8

2

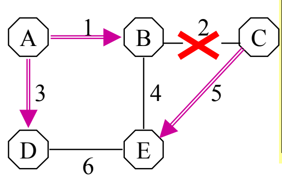
B generates a new DV
B=0,A=1,D=2,C=9,E=1

⇒ Distance seen by A to C grows to 10

A sends a new DV

1

A=0,B=1,D=1,C=10,E=2



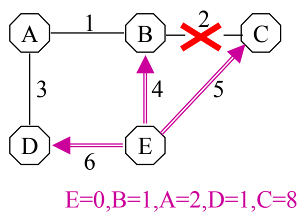
x to C	Link from x	Cost
A->C	1	10
B->C	1	11
C->C	-	0
D->C	3	11
E->C	5	8

2

B generates a new DV
B=0,A=1,D=2,C=11,E=1

E generates a new DV
E=0,B=1,A=2,D=1,C=8

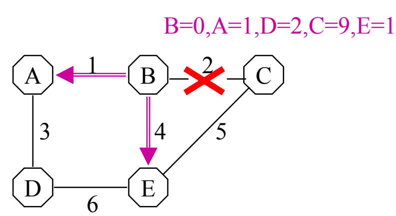
E sends a new DV



x to C	Link from x	Cost
A->C	1	10
B->C	4	9
C->C	-	0
D->C	6	9
E->C	5	8

B send its DV but the Tables are already OK

x to C	Link from x	Cost
A->C	1	10
B->C	4	9
C->C	-	0
D->C	6	9
E->C	5	8

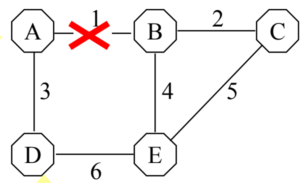


- Each update round improved the costs by 2
- The Process progresses in a random order, because it is genuinely parallel in nature.
- During the process, the state of the network is bad. DV-packets may be lost due to the overload created by bouncing user messages

Counting to Infinity occurs when failures beak the network to isolated islands (1)

- Linkki 1 is broken, and the network has recovered.
- All link costs = 1

A to	Link	Cost
D	3	1
A	-	0
B	3	3
E	3	2
C	3	3

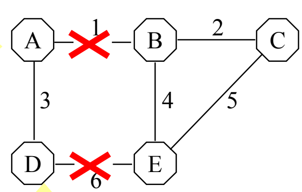


D to	Link	Cost
D	-	0
A	3	1
B	6	2
E	6	1
C	6	2

Counting to Infinity occurs when failures beak the network to isolated islands (2)

- Also link 6 breaks.
- D has not yet sent its distance vector.

A to	Link	Cost
D	3	1
A	-	0
B	3	3
E	3	2
C	3	3

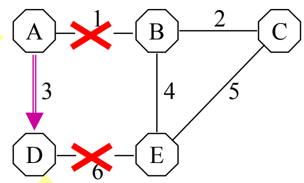


D to	Link	Cost
D	-	0
A	3	1
B	6	inf
E	6	inf
C	6	inf

Counting to Infinity occurs when failures beak the network to isolated islands (3)

- A sends its distance vector first.
  A=0,B=3,D=1,C=3,E=2
- D adds the information sent by A into its routing table.

A to	Link	Cost
D	3	1
A	-	0
B	3	3
E	3	2
C	3	3

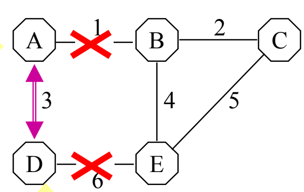


D to	Link	Cost
D	-	0
A	3	1
B	3	4
E	3	3
C	3	4

Counting to Infinity occurs when failures beak the network to isolated islands (4)

- The result is a loop. Costs are incremented by 2 on each round.
- An agreement is needed: Cost greater than any route cost is = inf.

A to	Link	Cost
D	3	1
A	-	0
B	3	5
E	3	4
C	3	5



D to	Link	Cost
D	-	0
A	3	1
B	3	4
E	3	3
C	3	4

Loops can be often avoided if less info is sent and by generating DVs immediately on RT change

Split horizon rule:

If node A sends to node X thru node B, it does not make sense for B to try to reach X thru A

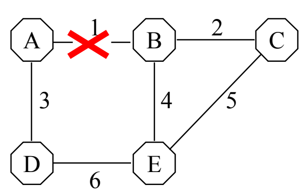
⇒ A should not advertise to B its short distance to X

Implementation choices:

1. A does not advertise its distance to X towards B at all
   ⇒ the loop of previous example can not occur
2. A advertises to B: X=inf. ("split horizon with poisonous reverse")
   ⇒ two node loops are killed

Three node loops are still possible (1)

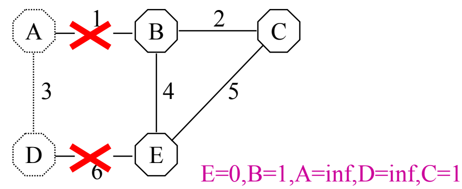
- Linkki 1 is broken, and the network has recovered.
- All link costs = 1



x to D	Link from x	Cost
B→D	4	2
C→D	5	2
E→D	6	1

Three node loops are still possible (2)

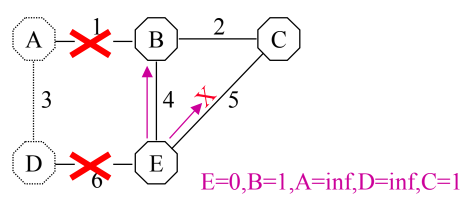
- Also link 6 fails.
- E sends its distance vector to B and C
  E=0,B=1,A=inf,D=inf,C=1



x to D	Link	Cost
B→D	4	2
C→D	5	2
E→D	6	inf

Three node loops are still possible (3)

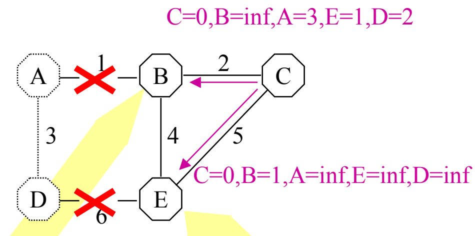
- Also link 6 fails.
- E sends its distance vector to B and C
  E=0,B=1,A=inf,D=inf,C=1
- ... But the DV sent to C is lost



x to D	Link from x	Cost
B→D	4	inf
C→D	5	2
E→D	6	inf

Three node loops are still possible (4)

- Now C sends its poisoned DV

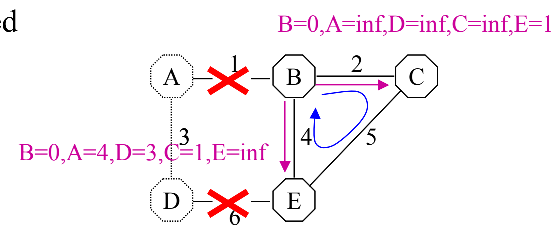


B to	Link	Cost
B	-	0
A	2	4
D	2	3
C	2	1
E	4	1

E to	Link	Cost
B	4	1
A	6	inf
D	6	inf
C	5	1
E	-	0

Three node loops are still possible (5)

- B generates its poisoned distance vectors
- The three node loop is ready
- Routes to D do not change except that the costs keep growing, nodes count to infinity.This finally breaks the loop: on link 5 cost=4 is advertised. C's knowledge about the distance to D grows ...



x to D	Link from x	Cost
B→D	2	3
C→D	5	2
E→D	4	4

When should a DV-protocol advertise

Time of advertisement is a compromise:

- (+) immediate delivery of change info
- (+) recovery from packet loss
- (+) need to monitor the neighbors
- (-) sending all changes at the same time
- (-) traffic load created by the protocol

(+) = Faster

(-) = Slower

Event triggered updates improve the functioning of RIP

- Entries in the Routing Tables have refresh and obsolescence timeouts.
- RIP advertises when the refresh timer expires and when a change occurs in an entry.
- Triggered updates speed up counting to infinity and reduce the probability of loops

The Bellman-Ford algorithm

Bellman-Ford algorithm (1)

- DV-protocols are based on the Bellman-Ford algorithm
- Centralized version:
  1. Let N be the nrof nodes and M the nrof links.
  2. L is the link table with M rows, L[*l*].m - link cost,
     L[*l*].s - link source
     L[*l*].d - link sink
  3. D is N × N matrix, such that D[*i*,*j*] is the distance from *i* to *j*
  4. H on N × N matrix, such that H[*i*,*j*] is the link *i* uses to send to *j*

```
D | 1  ..  i  ..  N |
1 |                 |
: |                 |
j |    distance     |
  |     from i      |
  |      to j       |
: |                 |
N |                 |
```

Both direction are presented separately in the Link table!
A Column ≡ DV of the corresponding node

Bellman-Ford algorithm (2)

- Initialized Distance and Link matrices

D	1	..		..	N
1	**0**	∞	∞	∞	∞
:	∞	**0**	∞	∞	∞
	∞	∞	**0**	∞	∞
:	∞	∞	∞	**0**	∞
N	∞	∞	∞	∞	**0**

Distance matrix D

H	1	..		..	N
1	**-1**	-1	-1	-1	-1
:	-1	**-1**	-1	-1	-1
	-1	-1	**-1**	-1	-1
:	-1	-1	-1	**-1**	-1
N	-1	-1	-1	-1	**-1**

Link matrix H

***NB:** Link vector has both directions of a link separately.*
First in D-matrix appear one hop link distances, second two hop link distances etc.

Bellman-Ford algorithm (3)

1. Initialization: If $i=j$, then D[i,j] = 0, else D[i,j] = inf. Initialize ∀ H[i,j] = -1. *(previous slide)*
2. ∀ l and ∀ destinations k set i = L[l].s, j = L[l].d and calculate d = L[l].m + D[j,k]
3. If d < D[i,k], set D[i,k] = d; H[i,k] = l.
4. If at least one D[i,k] changed, GOTO 2, else END.

Bellman-Ford algorithm (4)

- Number of steps ≤ N
- Complexity O(M.N**2)

RIP protocol

RIP-protocol properties (1)

- Simple protocol. Used before standardization.
- RIP version 1 – RFC 1058.
- RIP is used inside an autonomous system
- RIP works both on shared media (Ethernet) and in point-to-point networks.
- RIP runs on top of UDP and IP.

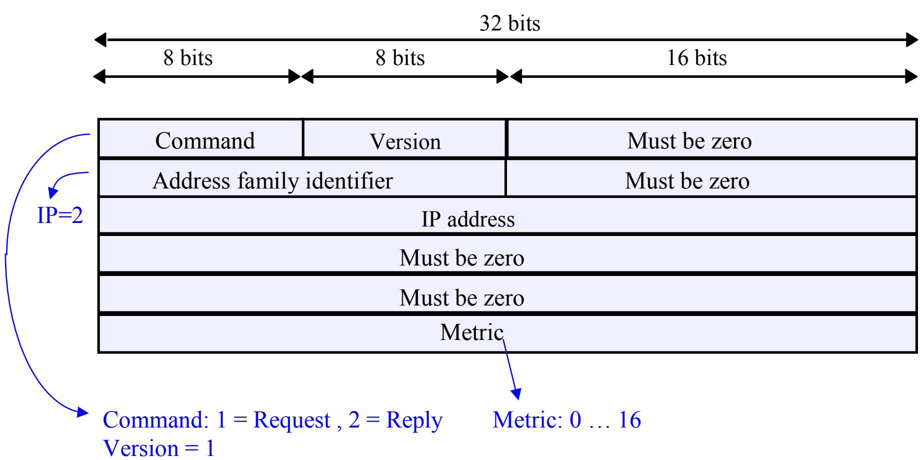
RIP-protocol properties (2)

- An Entry in the Routing Table represents a host, a network or a sub-net
  - <netid,subnetid,host> represents a host
  - <netid,subnetid,0> represents a sub-net
  - <netid,0,0> represents a network
  - <0.0.0.0> represents a route from the Autonomous System
- Information sent to the neighboring subnet is aggregated.

RIP-protocol properties (3)

- Distance = hop count = Nrof links on a path (route),
  - No other metrics
- RIP advertises once in 30s
  - if an entry is 180s old ⇒ distance is set to inf
- Distance 16 = infinite
- Timer triggered advertisements must be randomized to avoid bursts of RIP updates. 1-5 s.
- RIP uses poisoned vectors

RIP message format



RIP Routing Table

A routing table entry contains

- Destination IP address
- Distance to destination
- Next hop IP address
- "Recently" updated flag
- Several timers (refresh, obsolescence...)

RIP reply messages

- Distance vectors are sent in reply messages
- 30 seconds period
  - All routing table entries
  - Different DV on different links because of poisoned vectors
- Update messages after changes
  - Changed entries
  - 1-5 seconds delay, so that the message contains all updates that are related to the same change
- Destinations with infinite distance can be omitted if the next hop is same as before.

RIP request messages

- The router can request routing tables from its neighbors at startup
  - Complete list
- Partial routing table
  - For debugging

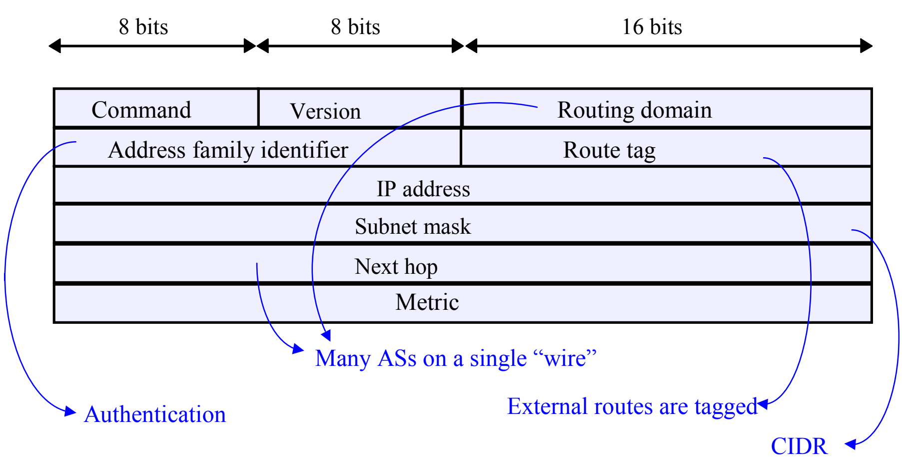
Silent nodes

- When only RIP was used, host could listen to routing traffic and maintain their own routing tables
  - Which router is closest to the destination?
  - Which link, if several available?
- These where ”silent nodes”, that only listened to routing traffic without sending
- Nowadays there are too many routing protocols
  - RIP-2, OSPF, IGRP, ...

RIP versio 2

- RFC-1388 (1387,1389)
- Why?
  - Simple and lightweight alternative to OSPF and IS-IS
- RIP-2 is a partially interoperable update with v1
  - RIP-1 router understand some of what a RIP-2 router is saying.
- Improvements
  - Authentication
  - Next hop –field
  - Subnet mask
  - External routes
  - Updates with multicast

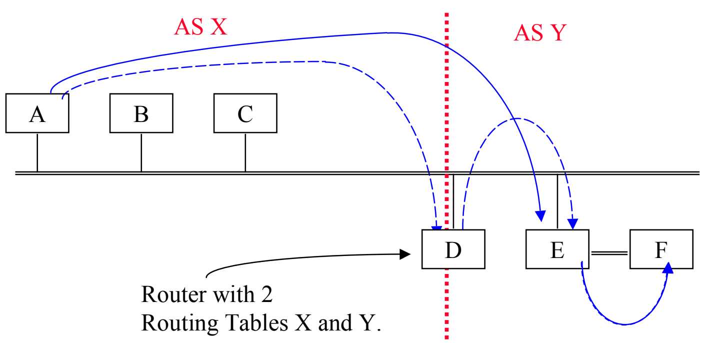
RIP versio 2 - sanomat



Routing from one sub-net to another

- In RIP-1 sub-net mask is not known outside the sub-net, only netid is sent in an advertisement out from a sub-net
  - ⇒ A host and a sub-net can not be distinguished
  - ⇒ All sub-nets must be interconnected with all other sub-nets and exterior traffic is received in the nearest router independent of the final destination inside our AS
- RIP-2 corrects the situation by advertising both the sub-net and the sub-net mask

Routing Domain and Next hop



Next hop ==> D advertises in X: the distance to F is f and the next hop is E!

Observations about RIP

- Routers have a spontaneous tendency to synchronize their send times. This increases the probability of losses in the net. Therefore, send instants are
  randomized between 15s ... 45s.
  - Reason: send interval = constant + time of message packing + processing time of messages that are in the queue.
- When RIP is used on ISDN links
  - A new call is established/30s ⇒ expensive.
- Slow network ⇒ queue length are restricted. RIP sends its DVs 25 entries/message in a row ⇒ RIP messages may be lost.
- A Correction proposal: ack all DVs: no periodic updates
  - ⇒ If there are no RIP message: assume that neighbor is alive and reachable
  - ⇒ Info on all alternative routes is stored.